

Avoiding Patterns in the Abelian Sense

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Abstract. We classify all 3 letter patterns that are avoidable in the abelian sense. A short list of four letter patterns for which abelian avoidance is undecided is given. Using a generalization of Zimin words we deduce some properties of ω -words avoiding these patterns.

1 Introduction

The study of words avoiding certain patterns is an area of combinatorics on words reaching back to at least the turn of the century and the work of Thue. A word in which no two adjacent subwords are identical is called *square-free*; thus a square-free word w can never be written $w = xyz$ with y a non-empty word. As an example, *orange* is square-free, but *banana* is not, since *banana* = *b an an a* contains *an* twice in a row.

One quickly checks that the longest square-free words over the alphabet $\{a, b\}$ are *aba* and *bab*; no word of length 4 or more over a 2 letter alphabet is square-free. Nevertheless, in 1906 Thue [17] showed that over a 3 letter alphabet, there are arbitrarily long square-free words. By König's Tree Lemma, this is equivalent to saying that there is an infinite square-free sequence on 3 letters. A word in which no three adjacent subwords are identical is called *cube-free*. Thue also showed that there is an infinite cube-free sequence over a 2 letter alphabet. A very readable modern account of Thue's work can be found in [3].

Thue's discoveries have numerous algebraic applications. One such is to the Burnside problem for groups. The solution of this problem by Novikov and Adjan [14] employed a cube-free sequence at a key point. More recently, square-free sequences have been used to study equational classes of semigroups [4] and non-finitely based varieties [16]. In this algebraic context, it becomes natural to consider not only square-free and cube-free words, but words avoiding other patterns.

Via a construction credited to Dilworth by Morse [13], the existence of an infinite square-free sequence on 3 symbols is equivalent to the existence of an infinite semigroup with 3 generators, obeying the law $x^2 = e$. Similarly, the existence of an infinite cube-free sequence on 2 symbols is equivalent to the existence of an infinite semigroup which obeys the law $x^3 = e$, and has 2 generators. In the study of semigroup varieties, the following Burnside type question is natural:

Given an arbitrary word p in one or more variables, is there a finitely generated infinite semigroup satisfying the law $p = e$?

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Pattern p	Smallest k such that p is k -avoidable
y	∞
y^2 (= yy)	3
y^n , $n \geq 3$	2

Table 1: The k -avoidability of patterns over a one letter alphabet.

Let Σ be a finite alphabet, and let Σ^* be the free monoid over Σ . Thus Σ^* consists of all finite strings of symbols in Σ , called *words*, with the operation of concatenation. The word in Σ^* with no letters, the *empty word*, is denoted ϵ . The length of word w is the number of letters in it, denoted by $|w|$. Thus $|banana| = 6$, $|\epsilon| = 0$. We will also consider infinite one-way words over Σ of the form $w = w_1w_2w_3 \dots$, which we call ω -words. The set of ω -words over Σ is denoted by Σ^ω . Words are finite unless otherwise indicated. If $p = p_1p_2 \dots p_k$ is a string and w is a finite or infinite word then w *encounters* p if there are words t, v, X_i such that $w = tX_1X_2 \dots X_kv$, each X_i is non-empty and $X_i = X_j$ whenever $p_i = p_j$. Equivalently, we say p *occurs* in w . We call p a *pattern* and write $w \vdash p$ if pattern p occurs in w . If p does not occur in w then we write $w \not\vdash p$ and say that w *avoids* p . Thus ‘banana’ encounters yy , while ‘orange’ avoids yy .

If arbitrarily long words avoiding p exist over an alphabet of size k , then p is *k-avoidable*, otherwise p is *k-unavoidable*. By König’s Tree Lemma, p is k -avoidable if and only if there is an ω -word on k letters that avoids p .

Thue’s work shows that yy and yyy are avoidable. In fact, everything about avoiding patterns on one letter has been known since 1906. (See Table 1.)

While some patterns are avoidable, others are unavoidable; consider xyx . If w is a long enough word over a fixed finite alphabet, then w has a repeated letter, say u . Word w will thus have a subword¹ uvu , some word v , and uvu is a coded version of xyx .

A characterization of avoidable patterns has been given by Zimin [18], and independently by Bean *et al.* [2]. According to Zimin, pattern p over an alphabet of n letters is avoidable if p is avoided by Z_n , where Z_n is defined recursively by

$$Z_n = \begin{cases} 1, & n = 1 \\ Z_{n-1}nZ_{n-1}, & n > 1. \end{cases}$$

The following problem has been open for 20 years [2]:

Find an algorithm to decide, given an avoidable pattern p , the smallest k such that p is k -avoidable.

Progress has been made on this problem given restrictions on p . For example, Cassaigne [6] solved the problem when p is a pattern over a 2 letter alphabet. Cassaigne also gives an inventory of the state of knowledge when p is a pattern over a

¹In this context, a subword of a word is always contiguous: *ana* is a subword (2 ways!) of *banana*, but *aa* is not. The opposite of a subword is an extension: *banana* is an extension of *ana*.

Pattern p	Smallest k such that p is k -avoidable in the abelian sense
y	∞
y^2 ($= yy$)	4
y^3	3
$y^n, n \geq 4$	2

Table 2: The k -avoidability in the abelian sense of y^n .

3 letter alphabet. For example, $abcbabc$ is known to be 3-avoidable, but perhaps is 2-avoidable.

In algebraic problems, commutativity is usually a simplifying assumption. In 1961, Erdős [10] asked whether an infinite sequence over a finite alphabet existed which was non-repetitive in an abelian sense; no 2 adjacent blocks were to be permutations of one another. Not only 1234 1234 is forbidden in such a sequence, but also 1234 2134, since 1234 and 2134 are permutations of each other.

A sequence avoiding yy in this abelian sense was discovered by Evdokimov [11], but he used an alphabet of 25 letters. The alphabet size was reduced to 5 letters by Pleasants [15], and not until 1992, to a 4 letter alphabet by Keränen [12]. One checks that on a 3 letter alphabet there are only finite sequences which are non-repetitive in this abelian sense. On the other hand, Dekking [9] showed that the smallest alphabet on which y^3 was avoidable in this abelian sense had 3 symbols, while 2 symbols were necessary and sufficient to avoid y^4 in the abelian sense. We see then, that our understanding of avoiding patterns in the abelian sense is only now at the same level that our understanding of ordinary pattern avoidance was in 1906!

If X and Y are words then we say that $X \sim Y$ if and only if the letters of X are a permutation of the letters of Y . In words, X is *abelian equivalent* to Y . If $p = p_1 p_2 \cdots p_k$ is a string and w is a finite or infinite word then w *encounters p in the abelian sense* if $w = tX_1 X_2 \cdots X_k v$, each X_i is non-empty and $X_i \sim X_j$ whenever $p_i = p_j$. Equivalently, we say p *occurs in w in the abelian sense*. We write $w \models p$ if w encounters p in the abelian sense, $w \not\models p$ otherwise. If an ω -word avoiding p in the abelian sense exists over an alphabet of size k , then p is *k -avoidable* in the abelian sense, otherwise p is *k -unavoidable* in the abelian sense.

From now on we have a distinction between ‘ w encounters p ’, and ‘ w encounters p in the abelian sense’, etc. In the first case we say by way of emphasis that ‘ w encounters p in the ordinary sense’. Suppose that P is a set of patterns. We extend our definitions, saying P is *k -avoidable* (*k -avoidable in the abelian sense*) if there is an ω -word w on a k -letter alphabet such that $w \not\models p$ ($w \not\models p$) for all $p \in P$. We write $w \not\models P$ ($w \not\models P$). Let S be a set of words. We say that w *avoids subwords in S* if no word of S is a subword of w .

Knowledge of abelian avoidability is in a very primitive state. It is not even known which patterns are abelian avoidable. We show (Lemma 3) that avoidability in the abelian sense is not the same as avoidability in the ordinary sense. In Section 3.1

we completely determine which patterns over a 3 letter alphabet are avoidable in the abelian sense.

Theorem 1 *A pattern p on 3 letters is avoidable in the abelian sense if and only if $Z_3 \not\models p$.*

To have the same result on 4 letters (see Section 3.2), it would suffice to show that the following six patterns are avoidable in the abelian sense.

- b_1 01020312
- b_2 01020321
- b_3 01021303
- b_4 01023031
- b_5 010203013
- b_6 010213020

For these 6 patterns there are considerable restrictions on an ω -word avoiding any of them. For example, none of these patterns is 4-avoidable in the abelian sense. If one of these patterns turns out to be avoidable in the abelian sense, then we have another contrast with ordinary avoidance, where no pattern is known to be 5-avoidable, but not 4-avoidable [1], [8].

The following conjecture seems natural:

Conjecture 2 *A pattern p on n letters is avoidable in the abelian sense if and only if $Z_n \not\models p$.*

2 Avoidable and Unavoidable Patterns

To begin with, we establish that avoidability in the abelian sense is not the same as avoidability in the ordinary sense.

Lemma 3 *The pattern $p = abcabadabacba$ is avoidable in the ordinary sense, but not avoidable in the abelian sense.*

Proof By Zimin’s criterion, to show that p is avoidable it is enough to show that p does not occur in $Z_4 = 121312141213121$ in the ordinary sense. Suppose on the contrary that Z_4 has a subword $P = ABCABADABACBA$ where A, B, C and D are non-empty words. Then $6|A| + 4|B| + 2|C| + |D| = |P| \leq |Z_4| = 15$, whence $|A| = |B| = 1$. By considering the frequencies of letters in P and Z_4 we see that $A = 1$ and $B = 2$. We see that the first and second occurrences of AB in P must correspond to the first and second occurrences of 12 in Z_4 , hence $C = 13$. On the other hand by considering the last and the next to last occurrences of BA in P , we deduce that $C = 31$, a contradiction. This shows that p is avoidable in the ordinary sense. On the other hand, the parsing $Z_4 = 1, 2, 13, 1, 2, 1, 4, 1, 2, 1, 31, 2, 1$ shows that p occurs in Z_4 in the abelian sense, which shows that p is unavoidable in the abelian sense, as Z_4 is unavoidable. ■

2.1 Avoidable Patterns

A *fractional abelian s power* is a word of the form uvw where $w \sim u$, and $|uvw|/|uv| = s$. For example, *madam* = *ma d am* is a fractional abelian $5/3$ power. We allow v to be empty in our definition, so it is correct to say that *teammate* = *team mate* is a fractional abelian 2 power, although we prefer to call it an abelian square. Also allowable is $u = w = \epsilon$, so it is also correct to say that *orange* is a fractional abelian 1 power. When u, v and w are all empty the ratio $|uvw|/|uv|$ is not defined, but it makes sense to say that $uvw = \epsilon$ is a fractional 0 power since $x^0 = \epsilon$ in the free monoid, and this is the convention we will make.

We say that a word w contains a r^+ *fractional abelian power* if w has a subword that is a fractional s abelian power with $s \geq r$. The notation $w \models x^{r+}$ indicates that w contains a fractional abelian r^+ power. Otherwise w *avoids* x^{r+} .

Theorem 4 ([7]) *Let $r \in (1, 2]$. There is an ω -word w over some finite alphabet such that w avoids x^{r+} .*

Call a word w *doubled* if each letter of w occurs at least twice in w .

Lemma 5 *Let p be a doubled pattern with $|p| = l$. Suppose $q \models p$. Then $q \models x^{r+}$ where $r = 1 + (l - 1)^{-1}$.*

Proof Let $p = a_1 a_2 \cdots a_l$ and suppose that $q = u A_1 A_2 \cdots A_l v$ where $A_i \neq \epsilon$ for each i and $A_i \sim A_j$ whenever $a_i = a_j$. Let $M = \max\{|A_k| : 1 \leq k \leq l\}$. Since p is doubled we may choose $i < j$ such that $a_i = a_j$ and $|A_j| = M$. Let $u = A_i, v = A_{i+1} \cdots A_{j-1}, w = A_j$. Then $u \sim w$ and

$$\begin{aligned} \frac{|uvw|}{|uv|} &= 1 + \frac{|w|}{|uv|} \\ &= 1 + \frac{|A_j|}{|A_i \cdots A_{j-1}|} \\ &\geq 1 + \frac{|A_j|}{(j-i)|A_j|} \\ &\geq 1 + \frac{1}{l-1}. \end{aligned}$$

Thus $q \models x^{r+}$, where $r = 1 + (l - 1)^{-1}$. ■

Lemma 6 *Any word w over an n letter alphabet with $|w| \geq 2^n$ contains a doubled subword.*

Proof Let $l_n, n \geq 1$ be the length of the longest word on n letters with no doubled subword. We shall prove by induction that $l_n \leq 2^n - 1$. Clearly $l_1 = 1 = 2^1 - 1$. Suppose the inequality holds for n and let w be a $n+1$ letter word. If w has no doubled subword then some letter, a say, occurs only once in w and $w = w_1 a w_2$ where w_1, w_2

are words on at most n letters. Neither of w_1 and w_2 can have a doubled subword since w has no doubled subword. By induction, $|w_1| \leq 2^n - 1$ and $|w_2| \leq 2^n - 1$. Thus $|w| \leq 2(2^n - 1) + 1 = 2^{n+1} - 1$, as desired. ■

Combining the last two lemmas and Theorem 4 gives us the following.

Lemma 7 Any pattern p on n letters with $|p| \geq 2^n$ is avoidable in the abelian sense.

We shall see that the avoidability of fractional abelian powers has further implications for the abelian avoidance of general patterns. Let p be a pattern and let $r \in [1, 2]$. We say that p forces x^{r+} if $w \models p$ implies $w \models x^{r+}$, and we write $p \rightsquigarrow x^{r+}$. If $p = \epsilon$ then p does not force any positive powers, since p may occur in the empty string and the empty string contains no positive powers (by our earlier convention). However, the empty string is just x^0 , so certainly $\epsilon \rightsquigarrow x^{0+}$. If $p \neq \epsilon$ then trivially $p \rightsquigarrow x^{1+}$, and hence also $p \rightsquigarrow x^{0+}$. Having dealt with the degenerate cases let us now show that $abacbab \rightsquigarrow x^{3/2+}$.

Lemma 8 The pattern $abacbab$ forces $x^{3/2+}$.

Proof Suppose $w = tA_1B_1A_2CB_2A_3B_3z$ where the A_i, B_i and C are non-empty words and $A_1 \sim A_2 \sim A_3, B_1 \sim B_2 \sim B_3$. If $|A_1| \geq |B_1|$, then $|A_1B_1A_2|/|A_1B_1| \geq 3/2$, so w contains the $3/2+$ power $A_1B_1A_2$. Otherwise, $|B_1| > |A_1|$, and w contains the $3/2+$ power $B_2A_3B_3$. ■

Let w be an ω -word avoiding $x^{3/2+}$ powers. Then w avoids $abacbab$. This method of showing avoidability can be generalized:

Let $p = p_1p_2 \cdots p_n$ be a pattern with letters p_i . Suppose that p is unavoidable in the abelian sense. Fix $r \in (1, 2]$, and let w be an ω -word avoiding x^{r+} . Since p is unavoidable, let $P = P_1P_2 \cdots P_n$ be a subword of w where the P_i are non-empty, and $P_i \sim P_j$ if $p_i = p_j$. Given r , and then P , define the P -length of a subword $T = p_i p_{i+1} \cdots p_j$ of p to be $|T|_P = |P_i| + |P_{i+1}| + \cdots + |P_j|$. Suppose that XYZ is a subword of p , with $X \sim Z$. Since w avoids x^{r+} , we must have $|XYZ|_P / |XY|_P < r$, or what is the same

$$(1) \quad |Z|_P < (r - 1)|XY|_P.$$

From the pattern p we then get a whole system of inequalities, one for each fractional power XYZ of p with $X \sim Z$.

Now let p be an arbitrary pattern, and let $r \in [1, 2]$. For each distinct letter p_j of p introduce a variable, l_j , which stands for the length of a string corresponding to letter p_j in an occurrence of p . For each subword $T = p_i p_{i+1} \cdots p_j$ of p we define l_T to be the sum $l_{p_i} + l_{p_{i+1}} + \cdots + l_{p_j}$. According to (1) above we create a system of inequalities $S(p, r)$ as follows:

$$S(p, r) : \{l_Z < (r - 1)l_{XY} : p = \alpha XYZ\beta, X \sim Z\}.$$

In the case where $r \in (1, 2]$ we exploit the fact that $l_X = l_Z$ and rewrite the system $S(p, r)$ in the equivalent form

$$S(p, r) : \{[(r - 1)^{-1} - 1]l_X < l_Y : p = \alpha XYZ\beta, X \sim Z\}.$$

Henceforth it is tacitly understood that if $r = 1$ the first form of the system is to be used, and the second, tidier form is to be used for all $r \in (1, 2]$.

Clearly, if $r = 1$ then system $S(p, r)$ is inconsistent, since the l_j are integer variables. In the case that p is unavoidable, we see that the system $S(p, r)$ is consistent for each $r \in (1, 2]$ by taking $l_{p_j} = |p_j|_p$ for each letter p_j of p as above. In the case of the simplest avoidable pattern, $p = aa$, and the largest exponent, $r = 2$, the system $S(aa, 2)$ defined above is just the inconsistency $0 < 0$, which shows that any ω -word avoiding x^{2+} powers automatically avoids p . More interesting is the case where $r \in (1, 2)$. For example, consider again Lemma 8, where $p = abacbab$ and $r = 3/2$. There $(r - 1)^{-1} - 1 = 1$ and the fractional abelian powers in p are aba , $abacba$, $bacb$, $bacbab$, $bacba$, $acba$, $abacbab$ and bab . The system $S(abacbab, 3/2)$ is

$$\begin{aligned} (2) \quad & l_a < l_b \\ (3) \quad & l_a < l_b + l_a + l_c + l_b \\ (4) \quad & l_a + l_b < l_a + l_c \\ (5) \quad & l_b < l_a + l_c \\ (6) \quad & l_b < l_a + l_c + l_b + l_a \\ (7) \quad & l_b + l_a < l_c + l_b \\ (8) \quad & l_b + l_a < l_c \\ (9) \quad & l_a < l_c + l_b \\ (10) \quad & l_a + l_b < l_a + l_c + l_b \\ (11) \quad & l_b < l_a. \end{aligned}$$

One immediately notices that the last inequality is the inversion of the first inequality. This is a contradiction! We deduce that $abacbab$ is avoidable, and in fact that $abacbab \rightsquigarrow x^{3/2+}$, in accordance with Lemma 8. When $S(p, r)$ is inconsistent we say that p has an *inversion*. All nonempty patterns have an inversion for $r = 1$, the *trivial* inversion. Typically when we speak of a pattern having an inversion we mean a nontrivial inversion with $r \in (1, 2]$. Note that the inequalities involving the variable l_c are of no use in deducing the inconsistency of $S(abacbab, 3/2)$. This is because variable l_c occurs only on the right hand side of the inequalities and it may be taken to be arbitrarily large. Indeed, for no value of $r \in (1, 2]$ will the presence of the variable l_c affect the consistency of the system, as we invite the reader to check. The fact that letter c occurs just once in p is what confines the singular appearances of variable l_c to the right hand sides of the inequalities. It makes sense to delete the inequalities involving l_c , and we should like to effect this deletion by deleting the letter c from $abacbab$. This motivates the discussion of the following paragraph.

Let $p_1 = a_1a_2 \cdots a_n$ and $p_2 = b_1b_2 \cdots b_m$ be patterns with occurrences $w = uA_1A_2 \cdots A_nv \models p_1$ and $w = sB_1B_2 \cdots B_mt \models p_2$. The occurrences of p_1 and p_2 are *consistent* if $A_i \sim B_j$ whenever $a_i = b_j$, and in that case we write $w \models p_1-p_2$. This leads to a generalization of the notion of a pattern, where we define an abelian occurrence of $p = p_1-p_2-\cdots-p_n$ in w to consist of occurrences of patterns p_1, p_2, \dots, p_n in w that are mutually consistent in pairs. We say that p_1, p_2, \dots, p_n are the n *components* of p , and that p is *connected* if $n = 1$ and *disconnected* if $n > 1$.

If $p = p_1-p_2-\cdots-p_n$ is a disconnected pattern and $r \in [1, 2]$ we define $S(p, r)$ to be $S(p_1, r) \cup S(p_2, r) \cup \cdots \cup S(p_n, r)$. The next lemma should come as no surprise.

Lemma 9 *Let $p \neq \epsilon$ be a connected or disconnected pattern, c a letter of p that occurs at most once in each component of p , and p' the pattern that results from p upon deleting letter c . Then for any $r \in (1, 2]$ the systems $S(p, r)$ and $S(p', r)$ are both consistent or both inconsistent.*

Proof Let the system $S_c(p, r)$ consist of those inequalities in $S(p, r)$ that involve the variable l_c . Then $S(p, r) = S_c(p, r) \cup S(p', r)$. Clearly if $S(p', r)$ is inconsistent then $S(p, r)$ is inconsistent. If on the other hand the system $S(p', r)$ is consistent then we see that the variable l_c occurs only on the right hand side of the inequalities of $S_c(p, r)$, and then exactly once in each inequality. This means that a given solution of the system $S(p', r)$ can be extended to a solution of $S(p, r)$ by taking l_c sufficiently large. This shows that the two system are both consistent or both inconsistent. The proof is complete. ■

Even if Lemma 9 cannot be applied there are still simplifications that can be made. Consider the substring $abacba$ of the pattern $p = abacbab$. The two parsings $a\ bac\ ba$ and $ab\ ac\ ba$ show that $abacba$ is both a fractional $6/5$ and a fractional abelian $3/2$ power. These two powers yield the inequalities (3) and (4) in $S(abacbab, 3/2)$ above. Clearly inequality (4) implies inequality (3), and this will still be the case if we vary $r \in (1, 2]$, so naturally we delete the weaker inequality. In general let $p = \alpha XYZ\beta = \alpha'X'Y'Z'\beta'$ with $X \sim Z$ and $X' \sim Z'$. For strings s and t we say that s is *lighter* than t if the multiset of letters of s is strictly contained in the multiset of letters of t , and we write $s \sqsubset t$. For example, $abac \sqsubset cbaadcad$, but $abac \not\sqsubset caadcad$. This gives a partial ordering of the substrings of p . We say that the fractional power XYZ *supersedes* the fractional power $X'Y'Z'$ if $X' \sqsubseteq X$ and $Y \sqsubseteq Y'$, and at least one of $X' \sqsubset X$ or $Y \sqsubset Y'$ holds. We write $X'Y'Z' \prec XYZ$ if XYZ supersedes $X'Y'Z'$. Finally, we prune $S(p, r)$ to the equivalent system

$$S(p, r)' : \{[(r - 1)^{-1} - 1]l_X < l_Y : p = \alpha XYZ\beta, X \sim Z, \text{ and } XYZ \text{ is } \prec\text{-maximal}\}.$$

For example, the system $S(abcba_cbabc, 3/2)'$ consists of just two inequalities, and is inconsistent. We define the *inversive threshold*, $r(p)$, of a pattern $p \neq \epsilon$ by the formula

$$(12) \quad r(p) = \inf\{r \in [1, 2] : S(p, r)' \text{ is inconsistent}\},$$

and by convention we set $r(\epsilon) = 0$. Note that $r(p)$ is defined for all patterns p , since $S(p, 1)$ is inconsistent for any $p \neq \epsilon$. Equivalently, for a nonempty pattern

$p = p_1-p_2-\dots-p_n$ we can define $r(p)$ as

$$(13) \quad r(p) = \inf_l \max \left\{ \frac{l_{XYZ}}{l_{XY}} : p_k = \alpha XYZ\beta, X \sim Z, 1 \leq k \leq n \right\},$$

where the infimum is taken over all functions l that assign positive integer lengths to the letters of p . We now describe an algorithm based on Lemma 9 that gives a lower bound for $r(p)$.

Given a pattern, p , we let $m(p)$ denote the multiset of letters of p , so $m(abacdabebaca) = \{a, a, a, a, a, b, b, b, c, c, d, e\}$. Given a function, l , that assigns positive integer lengths to the letters of p and two sub-multisets $U, V \subset m(w)$ we say that U is *longer than* V under l if the sum of the lengths of the letters in U is larger than the sum of the lengths of the letters in V . With $l: a \mapsto 1, b \mapsto 2, c \mapsto 3, d \mapsto 4, e \mapsto 5$ we have that $\{a, a, b, b\}$ is longer than $\{a, d\}$, but shorter than $\{a, c, c\}$. Given $p = \alpha XYZ\beta, X \sim Z$, and a letter j of X consider all assignments l with $l_j \geq l_i$ for all letters i in XYZ . In an occurrence of p where l gives the lengths of the strings corresponding to letters of p , what fractional power can we guarantee XYZ will correspond to? The answer is $(k + 2)/(k + 1)$, where $k \geq 1$ is the smallest integer for which there is a partition $m(Y) = Y_1 \cup Y_2 \cup \dots \cup Y_k$ such that $m(X)$ is longer than each $m(Y_i)$ under l . Specifically, since all we know is that l_j is largest among all l_i , this means that for $1 \leq m \leq k$ there are partitions $m(X) = J \cup O$ and $m(Y_i) = J' \cup O'$ such that J consists of all the letters j in X , $O' \subseteq O$ and $|J| \geq |J'|$. We write $R_j(X, Y, Z) = (k + 2)/(k + 1)$. For example, for p above we have $R_c(abac, dabe, baca) = 4/3$, since we have the partition $m(dabe) = \{a, b, d\} \cup \{e\}$, and no shorter partition will do. Similarly $R_a(abac, dabe, baca) = 4/3$ and $R_b(abac, dabe, baca) = 5/4$.

Given a non empty pattern $p = p_1-p_2-\dots-p_n$ and a letter j of p define $r_j(p)$ by

$$r_j(p) = \max\{R_j(X, Y, Z) : p_k = \alpha XYZ\beta, X \sim Z, 1 \leq k \leq n, \text{ and } j \text{ occurs in } X\}.$$

Thus $r_a(abacabc) = 3/2, r_b(abacabc) = 5/4$ and $r_c(abacabc) = 4/3$. Continuing with this example, suppose that $w \models abacabc$. If l_a is the largest among l_a, l_b, l_c then the string in w corresponding to aba is a $3/2+$ power. Similarly, if l_b is the largest then $bacab$ corresponds to a string in w that is a $5/4+$ power, and if l_c is largest then $cabcb$ corresponds to a $4/3+$ power. We conclude that $abacabc \rightsquigarrow x^{5/4+}$, as $5/4$ is the minimum of $3/2, 4/3$ and $5/4$. From this discussion it is clear that if we define

$$(14) \quad \underline{r}(p) = \min\{r_j(p) : j \text{ a letter of } p\},$$

then p forces a $\underline{r}(p)+$ power. If we define $\bar{r}(p)$ to be the largest fractional abelian power in p then this is clearly an upper bound for $r(p)$, since the infimum of equation 13 includes the function that assigns length 1 to each letter. We now have an algorithm for bounding $r(p)$ from below:

Theorem 10 Given a pattern $p = p_1-p_2-\dots-p_n$ the following algorithm returns a lower bound on the inversive threshold, $r(p)$:

- Step 0. If $p = \epsilon$ return(0).
- Step 1. While (\exists letter c occurring at most once in each component of p)
delete letter c from p .
- Step 2. If $p = \epsilon$ return(1),
else return($\underline{r}(p)$).

In particular, p has a nontrivial inversion if and only if Step 1 does not reduce p to the empty string.

Theorem 10 gives a combinatorial characterization of patterns that have inversions and it is also very easy to implement. That there are patterns p with $\underline{r}(p) < r(p) < \bar{r}(p)$ can be seen by taking $p = aba_cbc_bacb$, for which $\underline{r}(p) = 4/3$, $r(aba_cbc_bacb) = \sqrt{2}$ and $\bar{r}(p) = 3/2$. The determination of $r(p)$ is a fascinating open problem, but one that appears difficult.

2.2 Unavoidable Patterns

In seeking to show that a pattern p is unavoidable, one naturally investigates what kind of additional structure may be imposed on an ω -word avoiding p . An ω -word w is *recurrent* if every finite subword of w occurs infinitely often in w . If for each finite subword of a recurrent word w there is a bound on the distance between consecutive occurrences of that subword, then we say that w is recurrent *with bounded gaps*.

Theorem 11 *Let P be a set of patterns. If P is k -avoidable in the abelian sense then there is an ω -word on k letters that avoids P in the abelian sense and has the additional property of being recurrent with bounded gaps.*

Proof For definiteness, let $\Sigma = \{1, 2, 3, \dots, k\}$. Let

$$\mathcal{S} = \{S \subseteq \Sigma^* \mid \exists w \in \Sigma^\omega \ w \not\models P \text{ and } w \text{ avoids subwords in } S\}.$$

First notice that the set \mathcal{S} is not empty. Since P is k -avoidable, let $w \in \Sigma^\omega$ avoid P . Let q be a pattern in P . It follows that w avoids the word $1^{|q|}$ as a subword, since $1^{|q|} \models q$. Thus $\{1^{|q|}\} \in \mathcal{S}$.

We will show that the partial order (\mathcal{S}, \subseteq) is chain complete. Let $\{S_\alpha\}_{\alpha \in I}$ be a chain in the partial order, and put $S = \bigcup_\alpha S_\alpha$. We claim that $S \in \mathcal{S}$. To see this let $S_n \subseteq S$ be the set of all strings in S of length $\leq n$. Since S_n is finite and $\{S_\alpha\}_{\alpha \in I}$ is a chain, we have $S_n \subseteq S_\gamma$ for some $\gamma \in I$. Since $S_\gamma \in \mathcal{S}$ there is a $u \in \Sigma^\omega$ that avoids P in the abelian sense and avoids subwords in S_γ , hence u avoids subwords in S_n . The prefix of length n of u , u_n , also avoids P in the abelian sense, and furthermore u_n contains no string of S as a substring. Each word in the set $\{u_n\}_{n \geq 1} \subseteq \Sigma^*$ therefore avoids P and avoids subwords in S . By König's infinity lemma there is an ω -word $v \in \Sigma^\omega$ with both of these properties, that is v avoids P and avoids subwords in S . Therefore $S \in \mathcal{S}$, establishing our claim and proving that (\mathcal{S}, \subseteq) is chain complete. By Zorn's lemma, let $\hat{S} \in \mathcal{S}$ be a maximal element and let $w \in \Sigma^\omega$ avoid P in the abelian sense and avoid subwords in \hat{S} .

We claim that w is recurrent. Otherwise, some subword s occurs only finitely often in w . This means that some suffix w' of w does not contain s as a subword while avoiding P in the abelian sense. Then w' avoids subwords in $\hat{S} \cup \{s\} \in \mathcal{S}$. Since $\hat{S} \cup \{s\}$ is strictly larger than \hat{S} , we have a contradiction. Therefore, w avoids P in the abelian sense and is recurrent.

If some subword s of w has arbitrarily long gaps between its occurrences in w , then w has arbitrarily long subwords which avoid the subword s . By König's infinity lemma there is an ω -word w' on k letters that avoids P and avoids subwords in $\hat{S} \cup \{s\} \in \mathcal{S}$, again contradicting the maximality of \hat{S} . Therefore w is recurrent with bounded gaps. The proof is complete. ■

If w is a recurrent word and a pattern p occurs in w , then the substring of w that gives the occurrence of p is repeated farther out in w , so we have two identical occurrences of p in w separated by a non-empty substring. If 2 is some letter not in p corresponding to the separating substring then $p2p$ occurs in w . Again, $p2p3p2p$ occurs in w , where 3 is a letter different from 2 and not a letter of p , and so on. A definition is in order:

Definition 12 Let p be a word, none of whose letters is in \mathbf{N}^+ . The n -th Zimin word seeded with p is recursively defined by $Z_1(p) = p$ and $Z_n(p) = Z_n(p)(n)Z_n(p)$, $n \geq 2$.

Thus $Z_4(abc) = abc2abc3abc2abc4abc2abc3abc2abc$, for example. Of course the ordinary Zimin words $Z_n = Z_n(1)$, $n \geq 1$ are also seeded Zimin words. We now introduce functions ρ, λ, ι on words over the positive integers. Functions ρ and λ are morphisms generated by their action on integers:

$$\rho(i) = i + 1$$

$$\lambda(i) = \begin{cases} i - 1, & i > 1 \\ \epsilon, & i = 1. \end{cases}$$

Function ι is given by

$$\iota(w_1 w_2 w_3 \cdots w_n) = 1 w_1 1 w_2 \cdots 1 w_n 1,$$

where the w_i are positive integers.

Remark 13 Let π, κ be patterns over the positive integers with $\kappa = \kappa_1 \cdots \kappa_n$, $\pi = \pi' K_1 K_2 K_3 \cdots K_n \pi''$ where $K_i \sim K_j$ if $\kappa_i = \kappa_j$. We have

1. $\rho(\pi) \models \rho(\kappa)$
2. $\iota(\pi) \models \iota(\kappa)$ if neither of π and κ contains the letter 1
3. $\lambda(\pi) \models \lambda(\kappa)$ if $\kappa_i = 1$ whenever $K_i = 1$.

We also see that $Z_n = \iota\left(\rho\left(\lambda(Z_n)\right)\right)$ for all n . The following lemmas about Zimin words are readily established by induction:

Lemma 14 *If ABC is a subword of Z_n and $A \sim C$, then B contains some integer greater than any integer in A .*

Lemma 15 *If $Z_n \models p$ for some $n \geq 1$ then p is unavoidable in the abelian sense.*

We now give an application of Theorem 11 that generalizes Lemma 15 above. Suppose that $Z_n(p) \models q$ for some n and that p is unavoidable on k letters in the abelian sense. Then q is also unavoidable in the abelian sense on k letters. Otherwise by Theorem 11 there is a recurrent ω -word w on k letters that avoids q , and then $w \models Z_n(p) \models q$, a contradiction. If $Z_n(p) \models q$ for some n then we shall say that p generates q and write $p \rightarrow q$. Note that this relation is reflexive and transitive, but not symmetric. Let \mathcal{U}_k be the set of patterns on k letters that are unavoidable. We have proven the following:

Theorem 16 *If $p \in \mathcal{U}_r$ and $p \rightarrow q$, then $q \in \mathcal{U}_r$.*

An alternative, constructive way of looking at this result is the following: The polynomial $p_r(j) = \sum_{i=1}^j \binom{i+r-1}{r-1}$ counts the number, up to abelian equivalence, of non-empty words of length $\leq j$ on r letters. If $|w| = 2m + k$, $k \geq 0$ then there are $k(k+1)/2$ factorizations of w of the form $w = uX_1yX_2v$, where $|X_1| = |X_2| = m$, $y \neq \epsilon$, and u, v may be empty. The function $k(n) = \lceil \frac{1}{2}(-1 + \sqrt{8n+9}) \rceil$ then gives the smallest k such that a string of length $2m + k$ contains at least $n + 1$ disjoint, non-adjacent pairs of strings of length m . Now suppose that $p \rightarrow q$ and $p \in \mathcal{U}_r$. Then there is a minimum length ℓ_1 such that p occurs in any string on r letters of length ℓ_1 . We regard 2 occurrences of $p = a_1a_2 \cdots a_n$ to be the same up to abelian equivalence if for each i the block corresponding to a_i in the first occurrence is a permutation of the block corresponding to a_i in the second occurrence. Let n_1 be the number of ways p can occur in a string of length ℓ_1 , up to abelian equivalence. Any string on r letters of length $\ell_2 = 2\ell_1 + k(n_1)$ then contains an occurrence of $Z_2(p)$, and there are at most $n_2 = n_1 p_r(\ell_2 - 2|p|)$ different occurrences (up to abelian equivalence) of $Z_2(p)$ in such a string. In this way, by recursively defining $\ell_j = 2\ell_{j-1} + k(n_{j-1})$ and $n_j = n_{j-1} p_r(\ell_j - 2|Z_{j-1}(p)|)$ for $j \geq 2$, we see that $Z_t(p)$ occurs in any string on r letters of length ℓ_t . Since $p \rightarrow q$ we have $Z_n(p) \models q$ for some n , and it follows that q occurs in any string on r letters of length ℓ_n .

The next two lemmas are devoted to proving that if $p \rightarrow q$ then $Z_t(p) \models q$, where t is the number of distinct letters in q . Thus the problem of deciding if one pattern generates another one is decidable.

Lemma 17 *Suppose that q has t distinct letters and $Z_k \models q$ for some k . Then $Z_t \models q$.*

Proof Let t, k be as above, and take the letters of q to be positive integers. We will prove the theorem by induction on t . Suppose $t = 1$. Without loss of generality, $q = 1^r$ for some $r \geq 1$. However, $Z_k \models q$, and Z_k is free of abelian squares. It follows that $q = 1$, and $Z_1 \models q$. The theorem thus holds for $t = 1$. From now on, assume without loss of generality that both 1 and 2 are letters of q .

Now suppose that $t \geq 2$ and the theorem holds for $1, 2, \dots, t - 1$. We will use induction on k . In the case where $k \leq t$ we have $Z_t \models Z_k \models q$, and the theorem

holds. Suppose then that $k > t$. Fix an occurrence of $q = q_1q_2 \cdots q_n$ in Z_k ; say $Z_k = uR_1R_2 \cdots R_nv$ where $R_i \sim R_j$ whenever $q_i = q_j$.

If $\lambda(R_i) \neq \epsilon$ for each i then $\lambda(R_i) \sim \lambda(R_j)$ whenever $q_i = q_j$ and $Z_{k-1} = \lambda(Z_k) \models \lambda(R_1R_2 \cdots R_n) \models q$. By induction, $Z_t \models q$, and we are done.

Suppose then that $\lambda(R_{i_0}) = \epsilon$ for some i_0 , i.e., $R_{i_0} = 1$. We can suppose without loss of generality that $q_{i_0} = 1$. We now wish to use the special structure of Z_k . Every second letter of Z_k is a 1. If for some j we have $|R_j|$ even, then either we can write $R_j = 1Q_j$ for some word Q_j , or we can write $R_j = Q_j1$ for some word Q_j . Further, if $R_k \sim R_j$, then either $R_k = 1Q_k$ or $R_k = Q_k1$ where $Q_k \sim Q_j$. Analogous observations hold when $|R_j|$ is odd.

Form a new pattern $r = r_1r_2 \cdots r_n$ as follows:

$$r_j = \begin{cases} 1, & \text{if } R_j = 1 \\ 1q_j1, & \text{if } R_j = 1Q_j1 \text{ for some word } Q_j \neq \epsilon \\ 1q_j, & \text{if } R_j = 1Q_j \text{ for some word } Q_j \neq \epsilon \text{ not ending in a 1} \\ q_j1, & \text{if } R_j = Q_j1 \text{ for some word } Q_j \neq \epsilon \text{ not starting with 1} \\ q_j, & \text{if } R_j \text{ neither starts nor ends with a 1.} \end{cases}$$

We see that r has at most t distinct letters, one of which is a 1, since $R_{i_0} = 1$. By construction, $R_1R_2 \cdots R_n$ is an occurrence of r , and $r \models q$. By Remark 13, $Z_{k-1} = \lambda(Z_k) \models \lambda(r)$. Since 1 is a letter of r , $Z_{t-1} \models \lambda(r)$ by induction. It follows that $Z_t = \iota(\rho(Z_{t-1})) \models \iota(\rho\lambda(r)) \models q$, so we get $Z_t \models q$. This completes the proof by induction. ■

We can prove a more general version of the above lemma.

Theorem 18 *Suppose that q has t distinct letters and $Z_k(p) \models q$. Then $Z_t(p) \models q$.*

Proof For definiteness we assume that the letters of p and q are from disjoint alphabets and that these two alphabets contain no integers. We may assume that $k > 1$. Fix an occurrence of $q = q_1q_2 \cdots q_n$ in $Z_k(p)$; say $Z_k(p) = \alpha R_1R_2 \cdots R_n\beta$ where $R_i \sim R_j$ whenever $q_i = q_j$. Let δ be an operator that replaces substrings p by the letter 1. Thus, for example, $\delta(Z_k(p)) = Z_k$. It is possible that substrings p that overlap would render δ undefined, but we shall never apply δ in such a situation.

Case 1 Each R_i contains at least one integer. Let ϵ be an operator that deletes the letters of p . Then $Z_{k-1} \models \lambda\epsilon(Z_k(p)) = \lambda\epsilon(\alpha)\lambda\epsilon(R_1)\lambda\epsilon(R_2) \cdots \lambda\epsilon(R_n)\lambda\epsilon(\beta)$. Further, the $\lambda\epsilon(R_i)$ are non-empty, and $\lambda\epsilon(R_i) \sim \lambda\epsilon(R_j)$ whenever $q_i = q_j$. Thus $Z_{k-1} \models q$. By Lemma 17 we find that $Z_t \models q$, where t is the number of letters in q , hence $Z_t(p) \models Z_t \models q$, so $Z_t(p) \models q$, as desired.

Case 2 Each R_i contains only letters from p . In this case, $Z_t(p) \models p \models q$, hence $Z_t(p) \models q$, as desired.

Case 3 Some R_i contain only letters from p , but some R_i contain integers. In this second case, as we read from left to right across the pattern q we will encounter $n \geq 1$

letters x_1, x_2, \dots, x_n which correspond to those R_i in $Z_k(p)$ that contain integers. Factor q as $q = q_1 x_1 q_2 x_2 \cdots q_n x_n q_{n+1}$, and factor $Z_k(p)$ as $Z_k(p) = \alpha Q_1 X_1 Q_2 X_2 \cdots Q_n X_n Q_{n+1} \beta$ where $X_i \sim X_j$ whenever $x_i = x_j$ and the q_i are sub-patterns of q whose occurrences on the Q_i are compatible. None of the letters x_j appears in any q_i . Some of the q_i may be empty, but under our assumption at least one of them is not. We factor each string X_i as $X_i = u_i X'_i v_i$, $i = 1, 2, \dots, n$ where X'_i begins with the first integer in X_i and ends with the last integer in X_i . One or both of u_i, v_i may be empty, but if $x_i = x_j$ then $u_i v_i \sim u_j v_j$. Putting in the factorizations of the X_i we have

$$\begin{aligned} Z_k(p) &= \alpha Q_1(u_1 X'_1 v_1) Q_2(u_2 X'_2 v_2) \cdots Q_n(u_n X'_n v_n) Q_{n+1} \beta \\ &= \alpha' (v_0 Q_1 u_1) X'_1 (v_1 Q_2 u_2) X'_2 \cdots (v_{n-1} Q_n u_n) X'_n (v_n Q_{n+1} u_{n+1}) \beta' \\ &= \alpha' p X'_1 p X'_2 p \cdots p X'_n p \beta', \end{aligned}$$

where the strings v_0, u_{n+1} are borrowed from α and β respectively to complete the formation of the interspersing p 's.

The above factorization shows that $Z_k = \delta(Z_k(p)) \models 1x_1 1x_2 1 \cdots 1x_n 1 = q'$ (recall the definition of the δ operator after Theorem 18). By Lemma 17 we have $Z_{t'} \models q'$, where t' is the number of letters in q' .

Claim 19 *We can find an occurrence of q' in $Z_{t'}$, such that 1 in q' corresponds to 1 in $Z_{t'}$.*

Proof of Claim Write $Z_{t'}$ as $UA_1 Y_1 A_2 Y_2 \cdots A_n Y_n A_{n+1} V$ where $A_i \sim A_j$ for all i, j , while $Y_i \sim Y_j$ if $x_i = x_j$. By Lemma 14, each Y_i contains an integer higher than every integer in A_1 . If $A_1 = 1$, then we are done. Otherwise, none of the A_i or Y_i are 1, so by Remark 13,

$$\begin{aligned} Z_{t-1} &\models \lambda(Z_{t'}) \\ &= \lambda(UA_1 Y_1 A_2 Y_2 \cdots A_n Y_n A_{n+1} V) \\ &\models \lambda(A_1) \lambda(Y_1) \lambda(A_2) \lambda(Y_2) \cdots \lambda(A_n) \lambda(Y_n) \lambda(A_{n+1}) \\ &\models q' \end{aligned}$$

and the result can be established by induction. ■

Say then that $Z_{t'} = \alpha 1 Y_1 1 Y_2 1 \cdots 1 Y_n 1 \beta$ where $Y_i \sim Y_j$ whenever $x_i = x_j$. This yields an occurrence of q' in $Z_{t'}(p) = \alpha' p Y_1 p Y_2 p \cdots p Y_n p \beta'$ where 1 in q' corresponds to p in $Z_{t'}(p)$. Since $p = v_{i-1} Q_i u_i$ for $i = 1, 2, \dots, n + 1$ we have

$$\begin{aligned} Z_{t'}(p) &= \alpha' (v_0 Q_1 u_1) Y_1 (v_1 Q_2 u_2) X_2 (v_2 Q_3 u_3) \cdots (v_{n-1} Q_n u_n) Y_n (v_n Q_{n+1} u_{n+1}) \beta' \\ &= \alpha'' Q_1 (u_1 Y_1 v_1) Q_2 (u_2 X_2 v_2) Q_3 \cdots Q_n (u_n Y_n v_n) Q_{n+1} \beta'' \\ &= \alpha'' Q_1 Y'_1 Q_2 Y'_2 Q_3 \cdots Q_n Y'_n Q_{n+1} \beta'' \end{aligned}$$

where $Y'_i = u_i Y_i v_i$ for $i = 1, 2, \dots, n$. Since $u_i v_i \sim u_j v_j$ and $Y_i \sim Y_j$ whenever $x_i = x_j$ we have $Y'_i \sim Y'_j$ whenever $x_i = x_j$. Furthermore, for each i we may construct an occurrence of q_i on the string Q_i in such a way that all of these occurrences are compatible, simply by copying from the initial occurrence of q in $Z_k(p)$. We have now constructed an occurrence of q in $Z_{t'}(p)$.

Recall that at least one q_i is not empty, and the letters comprising the q_i are disjoint from the letters x_1, x_2, \dots, x_n . Thus in passing from $q = q_1 x_1 q_2 \cdots q_n x_n q_{n+1}$ to $x_1 x_2 \cdots x_n$ at least one letter is lost, and then by forming $q' = 1x_1 1x_2 1 \cdots 1x_n 1$ one letter is gained, so $t' \leq t$. Finally we have $Z_t(p) \models Z_{t'}(p) \models q$, hence $Z_t(p) \models q$, as desired. ■

Combining Theorem 16 with Theorem 18 we have:

Corollary 20 *If $p \in \mathcal{U}_r$ and $p \rightarrow q$, then q occurs in the abelian sense in any string on r letters of length ℓ_t , where ℓ_t is defined in the proof of Theorem 16 and t is the number of distinct letters in q .*

We now describe an idea that is useful in reducing the amount of backtracking needed to show that a pattern is unavoidable in the abelian sense on a given number of letters. If p and q do not contain the letter x then one may argue as in the proof of Theorem 16 that $p_q \in \mathcal{U}_k$ implies $pxq \in \mathcal{U}_k$, and hence both of these patterns require the same number of letters.

Lemma 21 *If p and q are patterns not containing the letter x then $p_q \in \mathcal{U}_r$ if and only if $pxq \in \mathcal{U}_r$.*

Proof If $pxq \in \mathcal{U}_r$ then clearly $p_q \in \mathcal{U}_r$. Suppose that $p_q \in \mathcal{U}_r$ but that $pxq \notin \mathcal{U}_r$. By Theorem 11 there is a recurrent ω -word w that avoids pxq . However, $p_q \in \mathcal{U}_r$ so p_q occurs in the abelian sense in w . The recurrence of w gives us two *identical* occurrences of p_q , which together yield an abelian occurrence of pxq , a contradiction. Therefore $p_q \in \mathcal{U}_r$ implies $pxq \in \mathcal{U}_r$ and the proof is complete. ■

To illustrate how Lemma 21 can be used consider the problem of showing that $p = 01021 \in \mathcal{U}_2$. The string 001100100 avoids p and has length 9, but $p_1 = 010_1$ occurs in any string of length 5 on 2 letters. Indeed, occurrences of p_1 are equivalent to occurrences of $p_2 = 010$, which is equivalent to $p_3 = 0_0$, which is the same as $p_4 = 0$, and the latter is certainly unavoidable.

3 Patterns on 3 and 4 Letters

3.1 The Case of 3 Letters

By Lemma 7, any pattern on 3 letters of length 8 or more is avoidable in the abelian sense. By [11] abelian squares are avoidable, hence so is any pattern encountering an abelian square. We need therefore only examine patterns of length at most 7 that are free of abelian squares. These are generated in Figure 3.1, where we assume that all patterns start with 01.

length	≤ 7	8	9	10	11	12	13	14	15	total
# of patterns	0	4	18	36	39	24	9	2	0	132

Table 3: Frequency distribution of the B-list.

b_1	01020312
b_2	01020321
b_3	01021303
b_4	01023031
b_5	010203013
b_6	010213020

Table 4: The six minimal patterns in the B-list.

patterns are abelian avoidable then it would follow that all patterns on 4 letters which do not encounter Z_4 are avoidable.

This last list of six we shall refer to as the *short B-list*.

4 Analysis of the Short B-List

In this section we will analyse the patterns of the short B-list and give some of their properties. The main outcome of our exploration is a lower bound on the number of letters needed to avoid these patterns and a further reduction of the B-list due to symmetries present in the patterns.

4.1 Symmetries

Let p^R denote the reverse of the pattern p , for example $abcbaac^R = caabcba$. If $w = pxq$, where p and q are words not containing the letter x then pxq^R is a *conjugate* of w . The word $w = abcadcabeabaa$ then has 2 conjugates: $w_1 = abcadaabaebac$ and $w_2 = abcadcabeaaba$. The word w is a *palindrome* if $w = w^R$.

Lemma 22 *If p, q do not contain the letter x and p is a palindrome then $f = pxq \in \mathcal{U}_k$ if and only if $\tilde{f} = pxq^R \in \mathcal{U}_k$.*

Proof By Lemma 21 $f \in \mathcal{U}_k$ if and only if $p_q \in \mathcal{U}_k$, which is equivalent to $q^R_p = q^R_p^R = (p_q)^R \in \mathcal{U}_k$ which is equivalent to $\tilde{f} = pxq^R \in \mathcal{U}_k$ by Lemma 21. ■

Note that if Conjecture 2 is true then it would follow that pxq is unavoidable if and only if pxq^R is unavoidable, p, x and q as in the hypotheses of Lemma 22. However, the statement of Lemma 22 is more specific since it involves k , the alphabet size. It is natural to make the following

Conjecture 23 *If p, q do not contain the letter x then $f = pxq \in \mathcal{U}_k$ if and only if $\bar{f} = pxq^R \in \mathcal{U}_k$.*

Given k , Lemma 22 says that the patterns 01021303 and 01023031 are both k -unavoidable or both k -avoidable in the abelian sense. This is also the case for 010203013 and 010213020, since 010213020^R is a conjugate of 010203013 with a palindromic prefix. This further reduces the short B -list to (say) the short, short B -list consisting of 01020312, 01020321, 01021303 and 010203013.

4.2 Minimal Lengths and Minimal Alphabets

In order to apply some of our lemmas to the patterns in the short, short B -list we need to find some patterns that generate these. We define the *spectrum of p* , denoted $Spec(p)$, as the set of all patterns that generate p . Since the property of one pattern generating another is transitive, we seek to find minimal elements of the spectrum of a pattern. Thus if $p \rightarrow q$ and $f \models p$ then $f \rightarrow q$, and we will not use f in our computations.

We find that $Z_4(00) \models 01021303, 010203013$. Also, $Z_3(00) \models 01020312, 01020321$. Thus 00 is in the spectra of these patterns. By the computation we did in Figure 3.1, all strings on 3 letters of length 8 or more contain an abelian square, so we look for shorter strings in the spectra. The only patterns in the tree of Figure 3.1 that do not occur in Z_3 are $g_1 = 0102101, g_2 = 0121012$ and $g_3 = 010212$. These therefore are the only patterns that can generate patterns on the short B -list. We find that 010212 does not generate any b_i , but that 0102101, 0121012 are in the spectra of all b_i . Also, no pattern on 4 letters of length 7 or less can generate any b_i , since these patterns either occur in Z_4 or encounter one of the g_i . We find that no patterns in $\{b_1, b_2, b_3, b_5\}$ can generate any other pattern in this set. We have

Lemma 24 *The patterns b_1 and b_2 occur in the abelian sense in any string of length 214 over a 3 letter alphabet, the patterns b_3 and b_5 occur in any string of length 208143 over a 3 letter alphabet.*

Proof Apply Lemma 20 with $r = 3, p = 00, \ell_1 = 8$, and $n_1 = 13$ to obtain $\ell_3 = 214$ and $\ell_4 = 208143$. ■

Lemma 25 *All patterns in the short B -list are unavoidable on 4 letters.*

Proof Utilizing a back tracking algorithm we find that the longest word on 4 letters that begins with 0102 and avoids in the abelian sense both 0102101 and 0121012 has length 54. Next we find that a longest string on 4 letters that begins with 0120 and avoids 0102101, 0121012 and *aba* as a substring has length 29. Finally we see that the longest string on 4 letters avoiding our patterns and with each four consecutive letters distinct has length 7. This proves that 0102101 and 0121012 are not simultaneously avoidable on 4 letters, hence all patterns in the short B -list are unavoidable on four letters since they each have 0102101 and 0121012 in their spectra. ■

This last result again underscores the difference between the abelian case and the ordinary case, as no pattern is known which is avoidable in the ordinary sense yet requires more than 4 letters to be avoided.

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